# Notes - Unit 2

# **OPTIMIZED IMPLEMENTATION OF LOGIC FUNCTIONS**

#### **BASIC TECHNIQUES:**

• We can always minimize logic functions using the Boolean theorems. However, more powerful methods such as Karnaugh maps and Quine-McCluskey algorithm exist: they provide a deterministic way to check that the minimal form of a Boolean function has been reached.

### KARNAUGH MAPS:

#### 2 variables:







#### 3 variables:









# ELECTRICAL AND COMPUTER ENGINEERING DEPARTMENT, OAKLAND UNIVERSITY ECE-378: Digital Logic and Microprocessor Design

#### Winter 2015

#### 4 variables:

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- This method appears in: "The Map Method for Synthesis of Combinational Logic Circuits", Maurice Karnaugh, *Transactions of the AIEE, Part I: Communication and Electronics*, vol. 72, no. 5, Nov. 1953, pp. 593-599. Karnaugh maps of 5, 6, 7, 8, and 9 are hinted at. Beyond 9 variables, the mental gymnastics for minimization are claimed to be formidable.
- The Quine-McCluskey algorithm provides a simpler approach when dealing with a relatively large number of variables.

## QUINE-MCCLUSKEY ALGORITHM:

- **Literal**: For an *n*-variable function *F*, it is a variable expressed as *X* or  $\overline{X}$ .
- Implicant: For an *n*-variable function, it is any product term that can appear in any possible sum of products (canonical or non-canonical) that represents the function. If P is an implicant, then P=1 implies that the function is 1. Thus, every minterm is an implicant.

A graphical way to see the implicants of a function is to take a look at the Karnaugh map (for a relatively low number of variables). All the possible terms we can get out of the K-map are implicants.

• Prime implicant: It is an implicant P such that the removal of any literal from P results in non-implicant of the function.

#### OUTLINE

1. Get the function to be minimized represented as a canonical Sum of Products: Use the minterm expansion form.

$$F(A, B, C, D) = \sum m(0, 1, 2, 5, 6, 7, 8, 9, 10, 14)$$

- 2. Get the Prime Implicants of the function: This is done by systematically applying  $XY + X\overline{Y} = X$  to all possible minterms and resulting non-canonical product terms. So, we build the Implicants Table by determining all Implicants:
  - ✓ We represent the minterms using the binary notation. For example:  $m_1 = \overline{AB}\overline{CD} = 0001$ . Then, we group the minterms by the number of ones they contain. For an *n* –variable function, the minterms have *n* literals.
  - Ve apply  $XY + X\overline{Y} = X$  to all possible pairs of minterms. This applies to pair of minterms that only vary by one literal. We attach a  $\checkmark$  to every minterm that was employed.

$$m_{0,1} = m_0 + m_1 = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D = \bar{A}\bar{B}\bar{C}$$

Note the table representation:  $m_{0,1} = m_0 + m_1 = 0000 + 0001 = 000 -$ . The symbol " - " indicates that a literal was simplified. The resulting column consists of terms with n - 1 literals.

✓ We keep applying  $XY + X\overline{Y} = X$  to all possible pair of resulting product terms. We attach a '✓' to every term that was employed. For each column we add, an extra literal is simplified (or a symbol " – " is added to the terms).

 $m_{0,1,8,9} = m_{0,1} + m_{8,9} = \bar{A}\bar{B}\bar{C} + A\bar{B}\bar{C} = \bar{B}\bar{C} \equiv 000 - +100 - = -00 - -0000 - -0000 - -000 - -000 - -000 - -000 - -000 - -00$ 

If we happen to get a repeated term, we eliminate one:

 $m_{0,1,8,9} = m_{0,8,1,9} = -00-$ ,  $\rightarrow m_{0,8,1,9}$  is eliminated

✓ When we cannot simplify any further, we stop and look for the terms that do not have a check '✓'. These terms are called the **Prime Implicants**. All the terms that appear in the table are the **Implicants**.

Number	4-literal	3-literal	2-literal	1-literal
of ones	implicants	implicants	implicants	implicants
0	$m_0 = 0000 \checkmark$	$\begin{array}{rcl} m_{0,1} &=& 000- \checkmark \\ m_{0,2} &=& 00-0 \checkmark \\ m_{0,8} &=& -000 \checkmark \end{array}$	$m_{0,1,8,9} = -00 - m_{0,2,8,10} = -0 - 0 - 0 - 0 - 0 - 0 - 0 - 0 - 0 - $	
1	m <sub>1</sub> = 0001 ✓ m <sub>2</sub> = 0010 ✓ m <sub>8</sub> = 1000 ✓	$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$	$m_{2,6,10,14} =10$ $m_{2,10,6,14} =10$	We can't combine any further, so we
2	$\begin{array}{cccc} m_5 = & 0101 & \checkmark \\ m_6 = & 0110 & \checkmark \\ m_9 = & 1001 & \checkmark \\ m_{10} = & 1010 & \checkmark \\ m_{m_{10}} = & 0111 & \checkmark \end{array}$	$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$		Stop Here
3	m <sub>7</sub> = 0111 ♥ m <sub>14</sub> = 1110 ♥			
4				]

 $F(A, B, C, D) = \overline{A}\overline{C}D + \overline{A}BD + \overline{A}BC + \overline{B}\overline{C} + \overline{B}\overline{D} + C\overline{D}$ 

- 3. Select a minimum set of Prime Implicants: *F* is the sum of this set that contains the minimum number of literals.
  - ✓ Build the Prime Implicant Chart. Mark the minterms that cover each single Prime Implicant with an 'X'.
  - ✓ Get the **Essential Prime Implicants**: Look for minterms that are covered by (are part of) a single Prime Implicant: this is, look for columns with one X. The corresponding Prime Implicants are the Essential Prime Implicants. The minimized *F* includes the Essential Prime Implicants. Thus, we must get rid of all the covered minterms of an Essential Prime Implicant: cross out the rows of the Essential Prime Implicants and the columns of the covered minterms. In the example, the Essential Prime Implicants are:  $\overline{BC}$ ,  $C\overline{D}$
  - ✓ For the remaining X's: select enough Prime Implicants to cover all the minterms of the function. This is a trial and error procedure: start by selecting the Prime Implicant that crosses out (rows and columns) most of the Xs, and so on.

Prime		Minterms										
Implicants		0	1	2	5	6	7	8	9	10	14	
<b>m</b> 0,1,8,9	ΒĒ	Х	Х					Х	x			
<b>m</b> 0,2,8,10	$\overline{B}\overline{D}$	Х		Х				Х		Х		
m <sub>2,6,10,14</sub>	CD			Х		Х				Х	x	
<b>m</b> 1,5	ĀĒD		Х		Х							
<b>m</b> 5,7	ĀBD				Х		Х					
<b>m</b> 6,7	ĀBC					Х	Х					

 $\rightarrow F(A, B, C, D) = \overline{B}\overline{C} + C\overline{D} + \overline{A}BD$ 

**EXAMPLE:**  $F(A, B, C, D) = \sum m(4, 8, 10, 11, 12, 15) + \sum d(9, 14)$ . Function with don't care terms.

✓ Implicants Table: To help simplifying the function, the don't care terms are included as minterms here. If a don't care term ends up being a Prime Implicant, we delete it (otherwise we are not taking advantage of the don't care terms).

Number	4-literal	3-literal	2-literal	1-literal
Number	4 IICEIAI	JIICEIAI	2 IICCIAI	I IICCIAI
of ones	implicants	implicants	implicants	implicants
0				
1	m <sub>4</sub> = 0100 ✓ m <sub>8</sub> = 1000 ✓	$\begin{array}{rrrr} m_{4,12} &=& -100 \\ m_{8,9} &=& 100 - \checkmark \\ m_{8,10} &=& 10 - 0 \checkmark \\ m_{8,12} &=& 1 - 00 \checkmark \end{array}$	$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$	
2	m <sub>9</sub> = 1001 ✓ m <sub>10</sub> = 1010 ✓ m <sub>12</sub> = 1100 ✓	$\begin{array}{c} m_{9,11} = 10 - 1 \checkmark \\ m_{10,11} = 101 - \checkmark \\ m_{10,14} = 1 - 10 \checkmark \\ m_{12,14} = 11 - 0 \checkmark \end{array}$	$m_{10,11,14,15} = 1-1-$ $m_{10,14,11,15} = 1-1-$	We can't combine any further, so we stop here
3	m <sub>11</sub> = 1011 ✓ m <sub>14</sub> = 1110 ✓	$m_{11,15} = 1-11 \checkmark$ $m_{14,15} = 111-\checkmark$		
4	m <sub>15</sub> = 1111 √			

F(A, B, C, D)	$= B\overline{C}\overline{D} + A\overline{D}$	$\overline{B} + A\overline{D} + AC$
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Prime Implicant Chart: The don't care terms are NOT included here. Only the minterms are included here, since we are trying to have as few X's as possible.

Prime		Minterms							
Implicants		4	8	10	11	12	15		
<b>m</b> 4,12	BŪD	x				Х			
<b>m</b> 8,9,10,11	$A\overline{B}$		х	Х	Х				
m <sub>8,10,12,14</sub>	$A\overline{D}$		Х	Х		Х			
<b>m</b> <sub>10,11,14,15</sub> AC				Х	Х		x		

• More than one minimal solution exist, depending on the **x** (in the same pink column) that we use:

 $\rightarrow F(A, B, C, D) = B\overline{C}\overline{D} + AC + A\overline{B}$  $Or: F(A, B, C, D) = B\overline{C}\overline{D} + AC + A\overline{D}$ 

**EXAMPLE**:  $F(A, B, C) = \sum m(0, 1, 2, 5, 6, 7)$ 

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Number		3-literal	2-literal	1-literal								
	of ones implicants		implicants	implicants								
	0	m₀ = 000 ✓	$\begin{array}{rcl} m_{0,1} &=& 00-\\ m_{0,2} &=& 0-0 \end{array}$									
	1	m <sub>1</sub> = 001 ✓ m <sub>2</sub> = 010 ✓	$m_{1,5} = -01$ $m_{2,6} = -10$	We can't combine any further, so								
	2	m <sub>5</sub> = 101 ✓ m <sub>6</sub> = 110 ✓	$m_{5,7} = 1-1$ $m_{6,7} = 11-$	we stop here								
	3	m7 = 111 ✓										

 $F(A, B, C) = \overline{A}\overline{B} + \overline{A}\overline{C} + \overline{B}C + B\overline{C} + AC + AB$ 

#### ✓ Prime Implicant Chart:

Prime		Minterms								
Implicants		0	1	2	5	6	7			
<b>m</b> 0,1	$\bar{A}\overline{B}$	х	Х							
<b>m</b> 0,2	ĀĒ	Х		Х						
<b>m</b> 1,5	ĒС		Х		Х					
<b>m</b> 2,6	ВĒ			х		Х				
<b>m</b> 5,7	AC				х		Х			
<b>m</b> 6,7	AB					Х	Х			

• No essential prime implicants. So, we can only select the minimum number of Prime Implicants that covers all the minterms. For the given arrangement, there is only one solution:

$$F(A, B, C) = \overline{A}\overline{B} + B\overline{C} + AC$$

- In the previous example, there were 2 solutions because we could pick any **x** in a column. Here, we can only pick one, that's why we have one solution.
- However, notice that there can be another way to cross out rows and columns producing a minimal solution:

Prime		Minterms								
Implicants		0	1	2	5	6	7			
<b>m</b> 0,1	$\overline{A}\overline{B}$	Х	Х							
<b>m</b> 0,2	ĀĒ	х		Х						
<b>m</b> 1,5	ĒС		х		Х					
<b>m</b> <sub>2,6</sub>	ВĒ			Х		Х				
<b>m</b> 5,7	AC				Х		Х			
<b>m</b> 6,7	AB					х	Х			

• Again, for this particular arrangement, there is only one minimal solution:  $F(A, B, C) = \overline{A}\overline{C} + \overline{B}C + AB$ 

#### ISSUES:

- To determine a minimal solution (i.e. solution with the same number of literals), we need to efficiently cross out rows and columns. We can do this by trial and error, but it can become a cumbersome procedure as the number of variables increase. And as illustrated in the example, there can be more than one way to efficiently cross out rows and columns.
- There can also be more than one minimal solution (even if there is only one way to efficiently cross out rows and columns) resulting from this method. We can determine all possible minimal solutions by inspection, but this can become cumbersome as the number of variables increase.
- A systematic way to determine all possible minimum solutions is provided by **Petrick's method**: given a prime implicant chart, we can determine all minimum sum-of-products solutions. This is out of the scope of this class.